Embeddings of Graphs and Hypergraphs in One–dimensional and Two–dimensional Grids

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A one dimensional embedding of an n node graph G(V, E) is a one-to-one mapping $\pi : V \to \{1, \ldots, n\}$. Three cost measures are usually associated with an embedding:

- 1. The bandwidth: $\max_{(u,v)\in E} |\pi(u) \pi(v)|$.
- 2. The cutwidth: $\max_{i=1}^{n-1} |cut(\pi^{-1}\{1,\ldots,i\},\pi^{-1}\{i+1,\ldots,n\})|.$
- 3. Sum of edge lengths: $\sum_{(u,v)\in E} |\pi(u) \pi(v)|$.

The problems of finding embeddings that minimize these cost measures are NP–Complete problems, and approximation algorithms have been proposed for them.

An overview of these approximation algorithms will be given. Generalizations to hypergraphs and two–dimensional embeddings will be discussed as well as related open problems.